Hamghilare : Natural Language Analysis System<br>TOKONAGA Takenobu, MWAYANA Mekoto, 'FANAKA Rownmi<br>Deparment of Comphter Science<br>Tokyo Institute ol Technology<br>KAMTWAKY Tadashi<br>Matachi Research Laboratory<br>Hitachi T,td.

## A Wentage

Thas paper perents a natual haguage amysis system Heugh A Th bused on gope Xe which passes with a bothomag and depthinse stratogy and has ability to lende loft aximposition. We have already developed a grammar for malism सest, which is a superset of DCG. With Xes , left extraposition phenoxema is natmally cemessed in gramom
 iments showed that in comparison to the oxiginal hill s. XGe systen, the analysis sped up 10 times in the interpeter mode and $4 x$ times in the compiled mode. The ThRE strucbued diehonary in hanglatis requires less memory, pro vides faster dickionary reterence and also handes complicated idioms with versatility. Consequently, the utilization of Hangite for practical purposes has becone frasible.

## 

So far, sevenal grammar fomalisin based on logic programming paradgm such os Metanomphosis Crammar [2] and DCG [9] have been presented. In Meanomphosis Crammar, cach grammar rule is translated into a Born Clause, and the Prolog interpreter parses the input sentence with these IIoma Clanse using a top-down and demb-fast strategy. Oulike ix the past where pasers had to be constructed for syutactic analysis, in this method, we do not have to becanse the Prolog interpreter itself worls as one. Metamorphosis Grammar also provides a natural language processing method which interleaves syntactic analysis und semantic analysis. This is a desirable feature from the poind of view of cognitive acience.

Following Metarorphosis Grammar, Pereira of al. developed a gronmar fomalism called Definite Clause Gammax(DCQ) and Bxamonibion Gammar(XG) [8]. The grammar rules written in $D C G$ are also translated into a Prolog program and the Prolog interpreter works as a topdown and depin first parser interleaving syatax analysis and semanhic analysis. XG is the extended vession of DCG capable of hauding lett cxiraposition.

However, top-down parser have a problem that the profram talls hato the infimite loop when a lefi recursive aule appears in the grammar rules. This problem can be solved by exher thansañg grammar rules with left recusive rales into ones whthout lett seciursive rules or by using a botiomup paring strategy Gince the former solution nay give omatnal yaving results, the latter is prefeable.

Maisumobo of Electrotechrical Laboratory developed is system in which the grammar roles written in DCG we tanmated robo Mom clauses called BOP clatses and Prolog intexpeter works as a botom-tip and deph-imed yarser


Figure 1: Stracture of LangLAB
for theso rules [14]. Matsumoto's system is called the BOP syatem, The BUY system can handle left recmave rules and, heat grammar mesi and the dictionary separately.

Komo of Tokyo Institute of Technology extended the BOP system to BOP- XG system [5] which cati handle the left extraposition phenomena elegandy. BUP-XG system introduced the grammar description form called X GS (Extraposition Grammar with Slash Category).

This paper preseats a natural languge analysis system Langlab based on Komo's BUP..XG system. Figure 1 shows the structure of the J,angLAB system. Users should poepare grammar rules written in XGS and a dictionary written in DCG. Both grammar rules and a dictionary are transluxed into BOP-XG clauses and TMRE structured dictionary respectively by translators. Translated results are consulted by the Prolog system and the Prolog interpreter works as a parser.

In chapter 2, we briefly explain the fundamentals of the BOP system and the grammar deseription form XGS adopted in LangLAB. We will also describe BOP-XG trantlator which traxslates the grammax writen in XGS into BUP XC clase and its optimizations. In chapter 3 , we will touch on the TRRE structured dictionay adoped in Langhab. The structured dictionary requires less memory and provides faster dicionary reference and provides flexible idiom handling. In chapter: 4, we shall present rosults of experiments verifying the effect of the optimization described in chapter 2. Experiments showed that the anal. ysis aped up 10 times in the interpretive mode and 4 times in the compiled mode. The authors believe that JangleAB performs well enough to be of practical use.

$$
\begin{array}{lll}
s & \rightarrow> & (d-1) \\
\text { np } & \rightarrow-> & \text { pron. } \\
\text { pron } \rightarrow->\text { [you]. } & (d-2) \\
\text { vp } \rightarrow \rightarrow \text { [日alk]. } & (d-3) \\
\hline
\end{array}
$$

Figure 2: Sample grammar written in DCG

```
np(G) --> {Link(np,G)}, (b~1)
    goal(vp),
    s(G).
pron(G) --> np(G). (b-2)
dict(pron) --.> [you]. (b-3)
dict(vp) --> [walk]. (b-4)
```

Figure 3: BOP clanses translated from figure 2

## 2 XGS and BUP-XG

In this chapter, we shall explain the grammar description form XGS adopted in LangLAB and the BUP-XG translator which translates grammar rules written in XGS into the BUP-XG clauses. Before explaining BUP-XG, we will briefly explain the mechanism of the BUP system, the predecessor of BUP-XG. Basic parsing mechanism of BUP is left-corner parsing with top-down prediction.

### 2.1 BUP system

In BUP system, grammar rules written in DCG (Figure 2) arc translated into the rules called BUP clauses which are also of DCG format, and some Prolog program (link clauses and termination clauses : explained later).

Figure 3 shows results of the translation. These BUP clanses are then translated into a Prolog program (ligure 4) by the DCG translator which is embedded in the Prolog system. Two more arguments are added to each predicate which denotes nonterminal symbol in figure 4. These arguments constitutes a difference list which represents the input string. With the special predicate goal which is necessary for bottom up parsing, this Prolog program can parse the input string with a bottom-up and depth-first strategy. Figure 5 shows the definition of the predicate goal.

Now, we shall give a step by step explanation of the parsing algorithm of the BUP system. We will use the grammar shown in figure 4 and input sentence "you walk" as an example. Calling the predicate goal activates the parsing process:

$$
\begin{aligned}
& \text { ?- goal(s, [you, malk], []). } \\
& n p(G, X, Z): \cdots \operatorname{link}(n p, G), \\
& \text { goal ( } \mathrm{vp}, \mathrm{X}, \mathrm{Y} \text { ), } \\
& \mathrm{s}(\mathrm{G}, \mathrm{Y}, \mathrm{Z}) \text {. } \\
& \operatorname{pron}(G, X, Y):-n p(G, X, Y) \text {. }(p-2) \\
& \operatorname{dict}(\text { pron, }[\text { you } \mid X], X) \text {. }(p-3) \\
& \text { dict ( } \mathrm{pp},[\text { ralk| } \mid \mathrm{X}], \mathrm{X} \text { ). }(\mathrm{p}-4)
\end{aligned}
$$

Figure 4: Prolog programs translated from figure 3

```
goal ( \(\mathrm{G}, \mathrm{X}, \mathrm{Y}\) ) :-
    \((\mathrm{g}-1)\)
    ( wf \(\quad\) goal \((G, X, \ldots)\)
        ;
        fail_goal \((G, X),!\), fail
        ),!
        wifgoal \((G, X, Y)\).
goal ( \(G, X, Y\) ) :- (g2)
        dict ( \(C, X, Y\) ),
        link( \(C, G\) ),
        \(P=\ldots[C, G, Y, Z]\),
        call(P),
        assertz (rf_goal(P)).
goal ( \(G, X, Y\) ) :- ( \(\mathrm{g}^{-3}\) )
        assertz(fail_goal(G,X)),!,
        fail.
```

Figure 5: Definition of the goal clause

This calling checks to see if :
A parse tree the root of which is the category " s ", can be constructed from the input string denoted by the difference between the list [you, wailk] and the list [ ] ([you, ralk] in this example).

The first call of goal invokes the clause ( $g-1$ ) in the figure 5 . The clause ( $\mathrm{g}-1$ ) checks to see if the same analysis have been made before, to avoid recompatation using the information previously asserted as wf_goal and fail_goal.

As the execution of the clause ( $g-1$ ) fails in this case, the system chooses the next clause ( $\mathrm{g}-2$ ). In the body of the clause ( $\mathrm{g}-2$ ), the system consults the dictionary by calling "dict ( $\mathrm{C},[$ [you, walk], $Y$ )". This predicate call picks ( $\mathrm{p}-3$ ) in figure 4 and the system matches "pron" with variable C and "[walk]" with variable $Y$.

In the second line of (g-2), the system calls the predicate link to see if the category which is obtained by the previous dictionary consultation ("pron" in this example) can be leftcorner of the current goal (" $s$ " in this example). The link clauses are calculated by the BUP translator. Suppose this test succeeds, the system calls the predicate "pron" :

```
P =.. [pron,s,[malk],[]],
call(P).
```

Calling "pron ( $s$, [ralk], [])" invokes ( $p-2$ ). Then, the system executes its body that is, "np(s, [walk], [])".

Calling "np ( $s$, [waik], [])" invokes the clause ( $p-1$ ). After calling the predicate link to check a reachability from "np" to "s", the system invokes "goal(vp, [walk], [])". At this point, the system has analyzed the string "you" as "np" and is predicting that the trailing string "walk" should be bundled up to the category " vp ".

In the same manner, a bottom-up analysis with a topdown prediction proceeds until the execution of goal with the termination clauses succeeds. See [14] for the detail of the termination clauses.

Results once succeeded or failed in an analysis are asserted as wf_goal in the end of ( $\mathrm{g}-2$ ) and fail_goal in the clause ( $g-3$ ) respectively. This information is used in ( $g-1$ ) as described.

| 8 | $\cdots>8$ rop, up. | ( $x-1$ ) |
| :---: | :---: | :---: |
| xp | $\cdots$ pron. | ( $x-2$ ) |
| xap | $\cdots{ }^{\text {c.i. }}$ det, nown, snel../np. | $(\mathrm{x}-3)$ |
| vp | $\cdots \gg \mathrm{vt}$, up. | ( $x-4$ ) |
| 3.rel | $\cdots$ xed pron, 3 。 | ( $x-5$ ) |

Figare 6: Sample grammar written in XGS

### 8.2 MUP $-X G$ system

The embedded semtence which appears in relative clauses in English can be viewed as a structure in which a noun phrase is missing from declarative sextence. A gap is formed as a result of moving the antecedent from within the declarative senteace to the left of the relative clause. Linguists call such phenomexta "Left extraposition". By considering the gap left by the movel constituents as a "trace", and incorporating a mechanism that looks for such a "trace" automatically, the number of grammar rules can be decreased and the graminar rules becom? casier to read. Moreover incorporating such mechanism contributes to making analysis speed faster.

Top-down parsers like ATNG [13], [12] and XG [8] incorporate such a mechanism. The top-down parser can predict what categocy the trailing input string may be bundled up to. Efficient trace searching is possible as the system assumes the existence of traces only when a particular category is predicted as a goal.

A pure boitom-up parser is not capable of such predictions and inefficieacy results because of the necessity to assume the existence of a trace between every two words. However, since the BUP system incorporates top-down prediction in the bottom-up parsing strategy described in 2.1, it is possible to implement the mechanism to look for the traces efficiently. Konno developed a BUP-XG system which incorporated such a mechanism [5].
The XGS adopted in LanghAB provides grammar writers the facility with which left extraposition can be naturally expressed in grammar rules. Figure 6 shows a small English grammar wlich is written in XGS.
The notation ". . /" (called "slash") in the rule ( $x-3$ ) is introduced in XGS. This rule means that there exists the syntactic category "np" which dominates the "trace" under the syntactic category "s_rel" ("s_rel" mean's relative sentence). This idea is influenced by the "slash category" in GPSG [3]. We call the category after "../" "slash category". Rule ( $x-3$ ) also shows that the category "np" consists of the categories "det", "noun" and "srel" and that the trace left behind by the left extraposition of the noun phrase consisting of "det" and "noun" is dominated by "s_rel". During an analysis, when the system finds the trace under "s.rel", as shown in figure 7, its associates the trace in the embedded sentence with the moved phrase ("the man").

XQS also provides a notation to represent "Ross's Complew NP constraint" [10]. Following is an example of this notation. "This notation is called "open ( $<$ )" and "close ( $>$ )" following Pereira [8].

$$
a \cdots b, c,\langle d\rangle .
$$

This rule mans that category "a" consists of categories " $b$ ", " $c$ " and " d ". Open-close notation defines the scope of extra-


Figure 7: Matching between slash category and its trace
position. This example says that the movement from under " b " or " c " to the outside of " a " is permissible, but the movement from under " d " to the outside of " a " is not. Sentences violating "Ross's Complex Np constraint" are rejected by modifying ( $x-3$ ) to become ( $x-3$ ')

$$
n p \quad \rightarrow>\text { det, noum, <s,rel../np>. }\left(x-3^{\prime}\right)
$$

With ( $x-3$ '), the trace which is dominated by slash category "np" under "s.rel" can only correspond to the noun phrase which consists of "det" and "noun".

In addition, XGS also provides a double arrow notation $(\Longrightarrow \gg)$ and the notation to describe X lists (explained later) explicitly. With these notation, "coordinate structure" can be represented in a natural way (see [5]).

### 2.3 BUP-XG translator

Just like the BUP system, the grammar rules written in XGS are translated into BUP-XG clauses, link clauses and termination clauses by the BUP-XG translator. The BUP$X G$ translator in the LangLAB system has been improved so as to generate BUP-XG clauses more optimized than that in the original BUP-XG system. Furthermore, it is also equipped with a new function which inserts parse tree information automatically. The translator takes about three seconds to translate a grammar of about 200 rules. The following subsection explains these improvements.

### 2.3.1 Representation of link clauses

As the number of grammar rules increases, more link clauses are generated by the translator. For example, from about 200 grammar rules of English which we have developed, the BUP-XG translator generates about 700 link clanses. Shortening the search time of link clauses would contribute to an efficient analysis.

Link clauses are called in the body of BUP-XCx clauses and in the predicate goal. Since both the arguments of link are atoms in the both cases, a link

## link (a,b).

which denotes the reachablity from the category "ib" to " b " can be change to the form

$$
a(b):=!
$$

This form of representation reduces the search space of the reachablity test. The BUP-XG translator in LangLAB generates link information of this form.

### 2.3.3 Hulexes foy diference list

As described in subsection 2.1 input string are represented by a difference list and intermediate analysis results are asserted with the predicate wrf goal and fail goal. Since the last two arguments of the wifgeal constitutes a differonce list of the input string, the longer the input string beeomes, the more memory wf goal consumes. By indexing difference lists, the amount of memory required is reduced, and faster reference to intermeriate results is possible.

For cxample, when the system gets the input string "you walk", the predicates text are asserted as follows:

```
text(s0, [1).
text(s1, [walk]).
text(s2, [you, walk]).
```

The dictionary reference program gets a difference list by calling text with indexes ( $s 1, s 2, \ldots$ ) as the key, before consulting the dictionary.

## 2. B. $^{3}$ Representation of intermediate results

Generally, a long input string gives rise to more wffgoals and fail goals which results in longer search time for intermediate analysis results. We-goals and fail-goals have as their arguments, the index to the difference list denoting the partial input string, and its analysis. As described in 2.1; goal first consults wf goals and fuil_goals with the indexes of input string as the key. In LangliAB system, the predicate names of intermediate analysis result are the indexes to the difference list instead of "wi_goal" or "fail_goal". This modification reduces the search space of the intermediato analysis results and speeds up the analysis process.

### 2.3.4 Insertion of parse tree information

Users sometimes require the results of syntactic analysis to be expressed as parse trees, and in both the BUP systern and the original BUP-XG system, users are required to inserf an argument in each category to accommodate parse tree information. However, it is not a difficult task to make the translator insert this information automatically. Yo the BUP-XG translator of JangLAB, this information is inserted automatically unless instructed otherwise. This function is similar to the one in the MicCord's MLG(Modular Logic Grammar) [7]. However, milike MLG, all the nonterminal symbols can be a node of parse trees.

### 2.3.5 Example of translation

Figure 8 shows the BUP-XG clauses translated from the grammar in figure 6. The variables beginning with " $X$ " in the figure 8 are introduced to handle leff extraposition. This variable is called X list (extraposition list) which were introduced in XG [8]. Infomnation pertaining to slash categories is pushed into the $X$ list and is then transferred from category to category during the analysis process. The predicate goal_x is an extended version of the predicate goal in the BUP system, which pops up the slash category from the X hist when the tace is found. Note that variables for parse tree information, the names of which begin with "T", are antomatically inserted and that the representation of link information (in braces) is also modified.

```
np (Goal, [TM], XMTO, X0, X1, Xh) \(\cdots \cdots\)
    [s(coal) \},
    goad x (vp,lT2l, Xi, xz),
    s(Goal, [Ts, TL, T2]], Into, X0, X2, XR).
```



```
    \(\{\operatorname{ng}(\) Goall \()\}\).
```




```
    \(\{\operatorname{mp}(G o a i)\}\),
    goal \(x\) (nown, [ [2] , \(\times 1, \times 2\) ),
    goal \(x(s, x e l,[r 31, x(\operatorname{np},[\operatorname{np}(t)], x 2), x 3)\),
    np (Goal, \([\) [np, T1, T2, T3] \(]\), Into \(, \times 0, X 3, X 10\) ).
vt (Goai, [T1], Thio, X0, Ki, XB) \(-\cdots\)
    \{ wp(Goal) \},
    goal \(x(\mathrm{mp},[\mathrm{T} 2], \mathrm{X}, \mathrm{X} 2)\),
    vp (Goa1, \([\) [wp, T1, T2], rmito \(X 0, X 2, x 6)\).
rel_prou(Goal, [ri], Thfo, X0, X1, XR) \(\cdots\)
    \{ s_rel(Goal) \},
    goal. x (s, \([\mathrm{ra}], \mathrm{Xi}, \mathrm{X} 2)\),
    \(s_{m}\) rel (Goal, \([\) [s, rel, \(\left.T 1, T 2]\right], \operatorname{Tnfo}, X 0, X Z, X 0\) ).
```

Figure 8: BUP-XG clauses translated from tigue 6

```
v(info(get)) --> [get].
v(Eef(get,[[vf|ed]])) >>> Lgot].
v(ref(get,[[vi|en]])) \cdots-s [gotten].
v(info(get,mp)) - -> [get, wp].
v(info(get on)) }->\mathrm{ [get, on].
```

Figure 9: Saxaple dictionary including idioms

## 3 THRTE structwed dictionary

This chapter explains the TRIE structured dictionary, another extension to the BUP-XG system and the BUP system. The TRLE strmetnred dictionary requires less monory, provides faster dictionary xeference and flexible idiom has. dling.

### 3.1 TRIW structure

The name "TRYE" is taken from "reTREEval" [1] and it means a kind of tree structure. A dictionary written in DCG is translated into a TRIE structured dictionary by the TRIE dictionary translator. The TRIE structure is a tapple which has three elements, that is "word", "information for word(s)" and "its child TRIE strecture".

For example, the dictionary written in DCG shown in figare 9 would be translated to the TRIE structured dictionary shown in figure 10 .

To look up a TRIE structured dictionary, the dictionary reference program searches through the tree matching the input string with the first element of the TRTE structure and, information for the string of input is retreved only after the last word of the input string is matched. Actually, the translator bundles up the dictionary entries which has, the same first word into a clause (see how the entries "get", "get on" and "get up" are translated in figure 10). By using this structure for the dictionary, the system can awoid the

```
4T%6|yeg.
    [Tv, Cumo(god)|]|.
    Clon,
                LIv, Tinco(p,t on)]I],
                11].
            mup,
                |v, Lino(gc&, mp)T]\,
            |] % )
dicha(got;
```



```
    [.].
dictalgotten,
```



```
    |!
```





La thace , the angument of the head is the infomation on
 of the catcy "क", Tho argumeat of the entry "got" amt "gothen" is a simetrae "ree" which denobes a pointer ro the cairy denoted by the first argument of "xas) (he thas case, a pontur to the entry "get"). Jictionary entries the fotomation of which only dillon from each other pextially, c.j. the root fome rad the congogeted fom of an inregular yob, cha be witton in this manner.

The second argument of the stracture "res" is the differental intommetion betwex this entry ("got" or "Foten") and the entsy pointed bo by the finet arganert of "pet ${ }^{5}$. Fa this oxaryphe, hanue "we" means "verb hom" and nits valao "od" and "en" means "yost" and "past partipho" wosec. Gively. Pith such a description, users do not have to winte addibional diom ont ses which melate the congugated form of magelar verbs. In the case of regolar verbs, wince coxjogabed forms ace processed by the homhologitel analysie progran buith io the dichonary relenemee program, idiom enbes whin include the conjugated fom are not nceessary. Tox axample, wsers do not have to mate the idiom entry "Kicked the breket", if the entry" kick the broker" is writton.

##  thouray


 in its "wout" position (firet elenext of trpplo). This teatre
 derents atech as "xot waly w bus also ...". Th the BUP" bysteve and the BUT TE system, ihe systom regads such idionas as bwo-dement word, that is a prefer texuinal part and a hombing nonterminal part. Tho fomer part is in-
 bine grammar roles. Yn Lraghab, the TRDE strectrod dichomacy is able to barde all mach idions as the dictomary Entres.

The idion entes whieh Hehude nu frozen elementa such m shown "not only a but mo wa be writen as fifure ll.

```
mg(D],LI) \cdots...% [not, onxy], mdj(\ldots..),
                lbot, also],adj(\ldots,.).
```



```
    [but,also], np(my2,\ldots);
    {jotu(ap1, Mp%, #p)}.
```

Nigun 11: Sample dictionary with nomtemmal symbols and mopeams in the xule body

```
diciabut, ll, !
    loney, \(\mathrm{A}, \mathrm{l}\)
        TLadj, .....1, 11, L
        Duterin. I:
            Talso, \(\square\), C
                Trady,\(\ldots\), ,
            [1ady, LD, []1J], [11]11]1],
    Cay, 1 , 1, I, I, 1
        [but.l], I
            [aiso, L], I
                LTmp, \(\mathrm{Hy} 2, \ldots, 1,1,1\)
                    (join(npl, wpa, wp)),
                    [Lup, (4p, (1111, [119119191111).
```



And figure $1 \%$ is the result of the translation.
m the case of BCC, as the idiom entry such as figme 11 is nomally handled as aganomar fulo, the momber of grammat rules increases and ineficioncy of analysis process results. It is preferable to hande gramomer rales and dictionary citries separately.

As shown in figure la, the translator converts the Prolog,

 orence program calls the proganm enclosed by parenthesis when it encounters such a fomm. In the same way, the synbuctic category in the dictionary entries such as "mp (Mpi,.)" are converted into the list the first element of which is a cabegory name and the rest of which are arguments of the cabegory (hap, Wp, , 7). The dictionary reference progam
 a form.

The TROE structured dictionary enables the IanghaB system to handle dioms with versatility [d].

## 4 Performance Considerations

We conducted experiments to vexify the effect of opimizar
 tic analysis of ten sample sentences. The experiment envj.roxmem is as lollows:

6 Machine : Sun $3 / 260$ Workstation

- Prolog: Quintus-Prolog Release 1.8

0 Grmmar: 16: rules in KCS
The the experiment, we measwed fhe time xequired to obhain all parse tree before and after the optimization for each

Table 1: Analysis time using interpretive code

| No. | Number of Words | Number of 'Trees | Analysis Time [msec] |  | $\begin{aligned} & \text { Ratio } \\ & (1) /(2) \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | (1) BUP-XG | (2) LangLAB |  |
| 1 | 14 | 9 | 80,4]5 | 8,552 | 9.40 |
| 2 | 4 | 12 | 18,868 | 2,700 | 6.99 |
| 3 | 3 | 7 | 46,700 | 4,983 | 9.37 |
| 4 | 1 | 10 | 30,900 | 3,600 | 8.58 |
| 5 | 3 | 11 | 39,634 | 4,050 | 9.79 |
| 6 | 4 | 18 | 95,933 | 9,550 | 10.05 |
| 7 | 9 | 21 | 323,167 | 26,183 | 12.34 |
| 8 | 2 | 19 | 87,550 | 9,349 | 9.36 |
| 9 | 4 | 17 | 180,300 | 15,816 | 11.40 |
| 10 | 1 | 25 | 116,284 | 12,083 | 9.62 |
|  |  |  |  | average | 9.69 |

## Table 2: Analysis time using compiled code

| No. | Number of Words | Number <br> of Trees | Analysis Time [mec] |  | $\begin{aligned} & \text { Ratio } \\ & (1) /(2) \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | (1) BUP-XG | (2) LangLAB |  |
| 1 | 14 | 9 | 20,485 | 4,134 | 4.96 |
| 2 | 4 | 12 | 2,467 | 1,299 | 1.90 |
| 3 | 3 | 7 | 4,783 | 2,284 | 2.09 |
| 4 | 1 | 10 | 2,884 | 1,566 | 1.84 |
| 5 | 3 | 11 | 4,383 | 1,917 | 2.29 |
| 6 | 4 | 18 | 18,768 | 4,500 | 4.17 |
| 7 | 9 | 21 | 127,400 | 14,000 | 9.10 |
| 8 | 2 | 19 | 13,450 | 4,450 | 3.02 |
| 9 | 4 | 17 | 59,468 | 8,2.16 | 7.24 |
| 10 | 1 | 25 | 23,650 | 5,801 | 4.08 |
|  |  |  |  | average | 4.07 |

sample sentence. 'This analysis does not include morphological analysis. Table 1 is the result of the experiment in the interpretive mode and table 2 is the one in the compiled mode. The fourth and the fifth column of the table is the time to analyze the sentence in the original BUP-XG system and in the LangLAB system respectively. Time is shown in millisecond.
Results showed that in comparison to the original BUP$X G$ system, the analysis sped up 10 times in the interpretive mode and 4 times in the compiled mode. The optimization is less effective in the compiled mode than in the interpretive mode. However, this optimization is practical because debugging is usually done in the interpretive mode. We believe that LangLAB has the capacity for practical use.
'There is a related work SAX [6] by Matsumoto. SAX is also a parsing system based on logic programming, but its parsing strategy is bottom-up and breadth-first. Okunishi of ICOT reports that LangLAB is $6 \sim 10$ times faster than SAX in the interpretive mode. However, in the compiled mode, SAX is $6 \sim 16$ times faster than LangLAB [11]. SAX has still yet to be modified to handle idioms. If this modification is introduced, debugging can be done on LangLAB in the interpretive mode and the debugged grammar can be executed on SAX in the compiled mode.

## 5 Conchusion

We have made the following modification to the original BOP-XG:

- Optimized and enhanced translated code
- Adopted TRIE structured dictionary

With these modifications, the analysis sped up in comparison to the original BUP-XG system and flexible idiom handling became possible. We believe that LangLAB has become a more powerful and practical tool for natural language processing. We plan to develop a natural language processing system which includes semantic analysis, based on LangLAB.

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